# **COMP6211I: Trustworthy Machine Learning Test-time Integrity (verification) part 2**

Minhao CHENG

Most slides are adapted from AAAI 2022 tutorial and from internet







## What is neural network verification?

#### Robustness



any noise in perturbation set

- Verification requires a formal proof to show the property holds •
- In the robustness verification setting, a model can't be attack  $\neq$  Verified lacksquare
- Many heuristic defense was broken under stronger attacks  $\bullet$
- A verified model cannot be attacked by any attacks (including unforeseen ones)





#### The Basic Formulation of Robustness Verification

Suppose  $f(x_0) > 0$ . Can we verify this property:

f(x) >



Must consider a set of infinite points as the input of the NN.

$$>0, orall x \in \mathcal{C}$$

Neural Network

#### The Basic Formulation of Robustness Verification

Assuming  $f(x_0) > 0$ , we solve the optimization problem to find the worst case:

 $\mathcal{C}$  is usually a perturbation set "around"  $x_0$ , e.g.,  $\mathcal{C} := \{x | \|x - x_0\|_p \le \epsilon\}$ 



Is it a hard problem?

 $f^* = \min_{x \in \mathcal{C}} f(x)$ 





# **CROWN** backward bound propagation

until we reach the input!

$$x W^{(1)} ReLU \hat{z}^{(1)}$$
  $ReLU f(x) \ge w^{(3) op} D^{(2)} W^{(2)} D^{(1)} W^{(1)} x + ext{constants}$ 

 $\min_{x} f(x)$ **CROWN linear bound:**  $x{\in}\mathcal{C}$ 

**computed efficiently** on GPUs in a backward manner

Goal: find linear relationships between output and every hidden neuron,



$$f(x) \geq \min_{x \in \mathcal{C}} oldsymbol{a}_{ ext{CROWN}}^ op x \! + \! c_{ ext{CROWN}} \! := \min_{x \in \mathcal{C}} f_{ ext{CROWN}}(x)$$

Where  $a_{\text{CROWN}}$  and  $c_{\text{CROWN}}$  are functions of NN weights, and can be

## The CROWN lower bound



#### Propagate bounds backwards



## $\alpha\text{-}\mathbf{CROWN}$ : further tighten the bounds

- ReLU neurons have a flexible lower bound for relaxation
- Try different lower bounds to find tightest bound
- Each unstable ReLU has a lower bound to select, so lots of freedom here



$$\hat{z}_j^{(i)} \ge \boldsymbol{\alpha}_j^{(i)} \boldsymbol{z}_j^{(i)} \ (0 \le \boldsymbol{\alpha}_j^{(i)} \le 1)$$

Adjustable!

 $\mathbf{l}_{i}^{(i)} \leq z_{i}^{(i)} \leq \mathbf{u}_{i}^{(i)}$  are pre-activation bounds, also computed using CROWN

## $\alpha$ -CROWN: further tighten the bounds

- Key idea: tighten bounds using gradients



 $f^*_{ ext{CROWN}} = \min_{x \in \mathcal{C}} oldsymbol{a}_{ ext{CROWN}}^ op x + c_{ ext{CROWN}}$ 

Actually a function of  $\alpha$ . How to effectively optimize  $\alpha$  to find the best bound?

$$\max_{lpha \leq 1} \min_{x \in \mathcal{C}} f_{ ext{CROWN}}(x;oldsymbol{lpha})$$

Inner minimization can usually be solved in closed form



# $\alpha\text{-}\mathbf{CROWN}$ : further tighten the bounds

- We can use gradient to optimize the relaxation, to make the bound tighter (tighter bound => stronger incomplete verification)
- We can make the bound **tighter** than the more expensive LP-based verifiers
- Optimization can be done rapidly on **GPUs**



#### **Branch and bound for ReLU Network Verfication**

Recall that ReLU neurons have three cases depending on pre-activation bounds:





# **Branch and bound: The branching step**

Split each "unstable" ReLU neurons to two subproblems:



Additional linear constraint ("split constraint"):

 $z_1$ 

 $z_1$ 

 $z_1 > 0$ OR  $z_1 < 0$ 

The additional constraint can make bounds tighter

# **Branch and bound: The bounding step**

bound for each subproblem:



Using an incomplete solver (traditionally, LP-based verifier) to get the lower



### Branch and bound search tree

Branch and bound improves the lower bound

Lower bound = -3.0 (computed by a incomplete verifier)

Lower bound = min(-2.0, 0.5) = -2.0

Lower bound = min(-1.0, -0.5) = -1.0

Lower bound = -0.5

Lower bound = 0.1



## Branch and bound search tree

Branch and bound is complete if each relaxed subproblem (**with split constraint**) can be solved to optimal.

$$f^* = \min_{x \in \mathcal{C}} f(x)$$

f\* Complete Verification Branch and bound with split constraints





### Branch and bound search tree

#### Idea:

Combine rapid bound propagation based incomplete verifiers on GPUs with branch and bound (BaB) to achieve complete verification

#### **Outcome:**

up to 100-1000x faster than MIP based approach, enable us to scale complete verification to larger models





# constraints; CROWN *cannot* handle it





- To use branch and bound, bound propagation must incorporate the split
  - hidden neurons

Propagate linear bounds backwards





Deal with split constraints with Lagrangians 



**CROWN**:

 $\max_{eta \geq 0} \min_{x \in \mathcal{C}} w'$ **β-CROWN**:  $\min$ (x) $x \in \mathcal{C}, z_1^{(2)} < 0$ 

Lagrangian/KKT multipliers S is an diagonal matrix with +/-1 and 0

 $\min_{x\in\mathcal{C}}f(x)\geq\min_{x\in\mathcal{C}}w^{(3) op}D^{(2)}z^{(2)}+ ext{const.}$  Cannot handle split constraint

$$^{(3) op} D^{(2)} z^{(2)} + eta^ op S^{(2)} z^{(2)} + ext{const.}$$



Lagrangians are also propagated! 



Linear coefficients changed with one additional term during propagation

$$w^{(3) op} D^{(2)} z^{(2)} + ext{const.}$$

$$w^{(3) op} D^{(2)} + eta^ op S^{(2)} \Big) z^{(2)} + ext{const.}$$



#### **β-CROWN main theorem:** all split constraints

Compared to (vanilla) CROWN ( $\beta$ =0):



 $\min_{x \in \mathcal{C}, z \in \mathcal{Z}} f(x) \ge \max_{\beta \ge 0} \min_{x \in \mathcal{C}} (\boldsymbol{a} + \mathbf{P}\beta)^\top x + \mathbf{q}^\top \beta + c_{\mathbf{n}}$ 

$$\min_{x \in \mathcal{C}} f(x) \geq \min_{x \in \mathcal{C}} oldsymbol{a}^ op x + c$$

Different  $\beta$  corresponds to different bounds, and we can choose the tightest one





- Assume we have a base classifier f that maps inputs x to labels y, i.e., f(x) = y• The approach creates corrupted versions of the image x by applying Gaussian noise with 0
- mean and variance  $\sigma^2$ , i.e.,  $\eta \sim \mathcal{N}(0, \sigma^2 I)$
- Left figure: input sample x; Right figure: image corrupted with Gaussian noise  $x + \eta$ • A smoothed classifier g is obtained by outputting the majority vote of the prediction on many Gaussian-corrupted images  $x + \eta$ 
  - The added random noise improves the robustness to adversarial perturbations





- To design a smoothed classifier g at the input sample x requires to identify the most likely class  $\hat{c}_A$ returned by the base classifier f on noisy images
  - Step 1: create *n* versions of x corrupted with Gaussian noise  $\eta \sim \mathcal{N}(0, \sigma^2 I)$
  - Step 2: evaluate the predictions by base classifier for all corrupted images,  $f(x + \eta)$
  - predictions for the second highest class  $\hat{c}_B$ ), return  $\hat{c}_A$  as the prediction by g(x)• Otherwise, if  $n_A - n_B < \alpha$ , abstain from making a prediction
  - Step 3: identify the top two classes  $\hat{c}_A$  and  $\hat{c}_B$  with the highest number of predictions on  $f(x + \eta)$ • Step 4: if  $n_A$  (number of predictions by f for the top class  $\hat{c}_A$ ) is much greater than  $n_B$  (number of





• Examples of noisy images from CIFAR-10 with varying levels of Gaussian noise  $\mathcal{N}(0, \sigma^2 I)$  from  $\sigma = 0$  to  $\sigma = 1$ 



 $\sigma=0.00$ 



 $\sigma = 0.25$ 



![](_page_23_Figure_7.jpeg)

- Intuitively, the certified radius R is large when:
  - The noise level  $\sigma$  is high
  - The probability of the top class
  - The probability of second top cl
- The authors prove that the certifie

- For binary classification,  $R = \sigma \Phi^{-1}(p_A)$ , because  $\Phi^{-1}(p_B) = -\Phi^{-1}(p_A)$

$$p_A = \mathbb{P}(g(x + \eta) = c_A)$$
 is high  
lass  $p_B = \mathbb{P}(g(x + \eta) = c_B)$  is low  
ed radius *R* is given by:

$$R = \frac{o}{2} \left( \Phi^{-1}(p_A) - \Phi^{-1}(p_B) \right)$$

•  $\Phi^{-1}$  is the inverse of the Gaussian cumulative distribution function

Certified robust radius by [Cohen et al.'19]:

![](_page_25_Figure_2.jpeg)

**Certified robust radius by [Cohen et al.'19]:** 

Classifier g(x) is certifiably correct for x, if

- 1. certified radius > adv budget
- 2. classifier g(x) is correct for x

Calculate the percentage of certifiably correct x and obtain certified accuracy for a dataset

![](_page_26_Figure_6.jpeg)

- Neyman-pearson
  - Given a sample from one of two distributions: null X or alternative Y
  - Two errors:
    - say "X" when the true answer is "Y"  $\rightarrow$  better
  - Optimal rule:

• say "Y" when the true answer is "X" -> limit its probability <= some failure rate  $\alpha$ 

• deterministically on the set  $S^* = \{z \in \mathbb{R}^d : \frac{\mu_Y(z)}{\mu_X(z)} \ge t\}$  for whichever t makes  $\mathbb{P}(X \in S^*) = \alpha$ .

![](_page_27_Figure_12.jpeg)

• Let  $X \sim \mathcal{N}(x, \sigma)$ 

By Lemma 3 it suffices to simply show that for any  $\beta$ , there is some t > 0 for which:

$$\{z:\delta^T z \le \beta\} = \left\{z:\frac{\mu_Y(z)}{\mu_X(z)} \le t\right\} \quad \text{and} \quad \{z:\delta^T z \ge \beta\} = \left\{z:\frac{\mu_Y(z)}{\mu_X(z)} \ge t\right\}$$

The likelihood ratio for this choice of X and Y turns out to be:

$$\frac{\mu_Y(z)}{\mu_X(z)} = \frac{\exp\left(-\frac{1}{2\sigma^2}\sum_{i=1}^d (z_i - (x_i + \delta_i))^2)\right)}{\exp\left(-\frac{1}{2\sigma^2}\sum_{i=1}^d (z_i - x_i)^2\right)}$$
$$= \exp\left(\frac{1}{2\sigma^2}\sum_{i=1}^d 2z_i\delta_i - \delta_i^2 - 2x_i\delta_i\right)$$
$$= \exp(a\delta^T z + b)$$

where a > 0 and b are constants w.r.t z, specificall Therefore, given any  $\beta$  we may take  $t = \exp(a\beta + b)$ , noticing that

$$\delta^T z \le \beta \iff \exp(a\delta^T z + b) \le t$$
$$\delta^T z \ge \beta \iff \exp(a\delta^T z + b) \ge t$$

$$\sigma^2 I$$
) and  $Y \sim \mathcal{N}(x + \delta, \sigma^2 I)$ .

ly 
$$a = \frac{1}{\sigma^2}$$
 and  $b = \frac{-(2\delta^T x + \|\delta\|^2)}{2\sigma^2}$ .

*Proof.* To show that  $g(x + \delta) = c_A$ , it follows from the definition of g that we need to show that

$$\mathbb{P}(f(x+\delta+\varepsilon)=c_A) > \max_{\substack{c_B \neq c_A}} \mathbb{P}(f(x+\delta+\varepsilon)=c_B)$$

We will prove that  $\mathbb{P}(f(x + \delta + \varepsilon) = c_A) > \mathbb{P}(f(x + \delta + \varepsilon) = c_B)$  for every class  $c_B \neq c_A$ . Fix one such class  $c_B$  without loss of generality.

For brevity, define the random variables

 $X := x + \varepsilon =$  $Y := x + \delta +$ 

In this notation, we know from (6) that

 $\mathbb{P}(f(X) = c_A) \ge \underline{p_A}$ 

$$= \mathcal{N}(x, \sigma^2 I)$$
$$\vdash \varepsilon = \mathcal{N}(x + \delta, \sigma^2 I)$$

and 
$$\mathbb{P}(f(X) = c_B) \leq \overline{p_B}$$

![](_page_29_Picture_11.jpeg)

![](_page_29_Picture_12.jpeg)

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and 
$$\mathbb{P}(f(X) = c_B) \leq \overline{p_B}$$

![](_page_30_Picture_11.jpeg)

![](_page_30_Picture_12.jpeg)

and our goal is to show that

 $\mathbb{P}(f(Y) = c_A)$ 

Define the half-spaces:

$$A := \{ z : \delta^T (z - x) \le \sigma \| \delta \| \Phi^{-1}(\underline{p}_A) \}$$
$$B := \{ z : \delta^T (z - x) \ge \sigma \| \delta \| \Phi^{-1}(1 - \overline{p}_B) \}$$

Algebra (deferred to the end) shows that  $\mathbb{P}(X \in A) = p_A$ . Therefore, by (8) we know that  $\mathbb{P}(f(X) = c_A) \ge \mathbb{P}(X \in A)$ . Hence we may apply Lemma 4 with  $h(z) := \mathbf{1}[f(z) = \overline{c_A}]$  to conclude:  $\mathbb{P}(f(Y) =$ 

Similarly, algebra shows that  $\mathbb{P}(X \in B) = \overline{p_B}$ . Therefore, by (8) we know that  $\mathbb{P}(f(X) = c_B) \leq \mathbb{P}(X \in B)$ . Hence we may apply Lemma 4 with  $h(z) := \mathbf{1}[f(z) = c_B]$  to conclude:  $\mathbb{P}(f(Y)) =$ 

of inequalities

$$\mathbb{P}(f(Y) = c_A) \ge \mathbb{P}(Y \in A) > \mathbb{P}(Y \in B) \ge \mathbb{P}(f(Y) = c_B)$$
(12)

$$_A) > \mathbb{P}(f(Y) = c_B) \tag{9}$$

$$c_A) \ge \mathbb{P}(Y \in A) \tag{10}$$

$$c_B) \le \mathbb{P}(Y \in B) \tag{11}$$

To guarantee (9), we see from (10, 11) that it suffices to show that  $\mathbb{P}(Y \in A) > \mathbb{P}(Y \in B)$ , as this step completes the chain

We can compute the following:

 $\mathbb{P}(Y \in A) = \Phi\left(\Phi^{-1}(Y \in B) = \Phi\left(\Phi^{-1}(Y \in B) = \Phi\left(\Phi^{-1}(Y \in B) = \Phi\left(\Phi^{-1}(Y \in B) \right)\right)\right)$ Finally, algebra shows that  $\mathbb{P}(Y \in A) > \mathbb{P}(Y \in B)$  if and only if:

 $\|\delta\| < \frac{\sigma}{2}(\Phi$ 

$$= \Phi\left(\Phi^{-1}(\underline{p}_{A}) - \frac{\|\delta\|}{\sigma}\right)$$
$$= \Phi\left(\Phi^{-1}(\overline{p}_{B}) + \frac{\|\delta\|}{\sigma}\right)$$

$$\Phi^{-1}(\underline{p_A}) - \Phi^{-1}(\overline{p_B}))$$

![](_page_32_Figure_6.jpeg)

![](_page_32_Figure_7.jpeg)

- Certified top-1 accuracy by ResNet50 on ImageNet with the random smoothing approach
  - Top row: the certified top-1 accuracy of 49% under adversarial perturbations  $\ell_2 < 0.5$ 
    - This is achieved with noise level  $\sigma = 0.25$
    - - Note that perturbation with  $\ell_2$  norm < 0.5 is fairly small
      - For example, perturbation with  $\ell_2 = 1$  can change one pixel by 1 (=255/255), or change 10 pixels by 0.3 ( $\approx 80/255$ ), or change 1,000 pixels by 0.03 ( $\approx 8/255$ )
  - Increasing the  $\ell_2$  radius from 0.5 to 3.0 reduces the certified accuracy
  - For comparison, the standard top-1 accuracy on clean images by the smoothed classifier g is 67%

$\ell_2$ radius	best $\sigma$	Cert. Acc (%)	STD. ACC(%)
0.5	0.25	49	67
1.0	0.50	37	57
2.0	0.50	19	57
3.0	1.00	12	44

• For any perturbation with radius  $\ell_2 < 0.5$ , the robust classifier will correctly predict the class